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**A CROWDED Q -POINT UNDER $\text{CPA}_{\text{prism}}^{\text{game}}$**

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ABSTRACT. In this note, we prove that the version $\text{CPA}_{\text{prism}}^{\text{game}}$ of the Covering Property Axiom, which holds in the iterated Sacks model, implies that there exists an ω_1 -generated crowded ultrafilter on \mathbb{Q} which is also a Q -point. Since no crowded ultrafilter can be a P -point, this constitutes an interesting example of a Q -point which is not a P -point.

1. INTRODUCTION

We will use standard set theoretic notation as in [7]. Let \mathcal{U} be a non-principal ultrafilter on a countable set X . Then, \mathcal{U} is a P -point if for every partition \mathcal{P} of X , either $\mathcal{U} \cap \mathcal{P} \neq \emptyset$ or there exists an $X \in \mathcal{U}$ such that $X \cap P$ is finite for each $P \in \mathcal{P}$. \mathcal{U} is a Q -point if for every partition \mathcal{P} of X into finite pieces, there exists an $X \in \mathcal{U}$ such that $|X \cap P| \leq 1$ for each $P \in \mathcal{P}$. Given a non-principal ultrafilter \mathcal{U} on X , we say that $\mathcal{B} \subset \mathcal{U}$ is a basis for \mathcal{U} if for every $U \in \mathcal{U}$ there exists a $B \in \mathcal{B}$ such that $B \subset U$. Then, we can define the *character* of \mathcal{U} as $\chi(\mathcal{U}) = \min\{|\mathcal{B}| : \mathcal{B} \text{ is a basis for } \mathcal{U}\}$. We say that \mathcal{U} is κ -generated if $\chi(\mathcal{U}) = \kappa$.

Consider \mathbb{Q} with the subspace topology induced by the usual topology on \mathbb{R} and denote by $\text{Perf}(\mathbb{Q})$ the family of its perfect

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subsets. A non-principal filter \mathcal{U} on \mathbb{Q} is *crowded* if the family $\text{Perf}(\mathbb{Q}) \cap \mathcal{U}$ forms a basis for \mathcal{U} . The crowded ultrafilters have been studied in connection with the remainder of the Stone-Čech compactification of \mathbb{Q} , and their existence follows from the Continuum Hypothesis, Martin's Axiom for countable posets [5], or from the equality $\mathfrak{b} = \mathfrak{c}$ [4].

In [2], K. Ciesielski and J. Pawlikowski showed that a version of their Covering Property Axiom called $\text{CPA}_{\text{prism}}^{\text{game}}$, which holds in the iterated Sacks model, implies that there exists an ω_1 -generated crowded ultrafilter on \mathbb{Q} , and they noted that no crowded ultrafilter can be a P -point. This result is interesting because CPA implies $\mathfrak{b} < \mathfrak{c}$.

The main result of this paper is that $\text{CPA}_{\text{prism}}^{\text{game}}$ implies the existence of an ω_1 -generated crowded ultrafilter on \mathbb{Q} which is also a Q -point¹. Notice that this contradicts the remark by Ciesielski and Pawlikowski in [2, p. 49] that crowded ultrafilters cannot be Q -points.

It is a result of Arnold W. Miller [9] that there are no Q -points in Richard Laver's model for Borel's Conjecture [8]. Since the equality $\mathfrak{b} = \mathfrak{c}$ holds in Laver's model, it is consistent with ZFC that no crowded ultrafilter on \mathbb{Q} is a Q -point.

2. PRELIMINARIES ON $\text{CPA}_{\text{cube}}^{\text{game}}$ AND $\text{CPA}_{\text{prism}}^{\text{game}}$

2.1 Cubes and Prisms

The framework of CPA rests on the concepts of *cube* and *prism*. If \mathfrak{C} denotes the space 2^ω with its usual product topology and \mathfrak{X} is a Polish space, then we define

$$\text{Perf}(\mathfrak{X}) = \{C \subset \mathfrak{X} : C \text{ is homeomorphic to } \mathfrak{C}\}.$$

A *perfect cube* in \mathfrak{C}^ω is any set $C = \prod_{i < \omega} C_i$ where $C_i \in \text{Perf}(\mathfrak{C})$ for every $i < \omega$. If \mathfrak{X} is a Polish space, then a *cube* in \mathfrak{X} is a pair $\langle f, P \rangle$ where $f: C \rightarrow \mathfrak{X}$ is a continuous injection and $P = f[C]$ for some perfect cube C . The following proposition is one of the principal tools for using CPA, and it is a refinement of a theorem proved independently by H. G. Eggleston [6] and M. L. Brodskii [1].

¹Recently the author has proven that $\text{CPA}_{\text{prism}}^{\text{game}}$ implies that there is also a crowded Q -point of character \mathfrak{c} .

Proposition 1 (Ciesielski and Pawlikowski [3, Claim 1.1.5]). *Consider \mathfrak{C}^ω with its usual topology and its usual product measure. If G is a Borel subset of \mathfrak{C}^ω , which is either of second category or of positive measure, then G contains a perfect cube.*

The notion of *prism* is a generalization of that of a cube. If $\alpha < \omega_1$ is a non-zero countable ordinal, let $\Phi_{\text{prism}}(\alpha)$ be the set of all functions $f: \mathfrak{C}^\alpha \rightarrow \mathfrak{C}^\alpha$ with the property that

$$f(x) \upharpoonright \xi = f(y) \upharpoonright \xi \Leftrightarrow x \upharpoonright \xi = y \upharpoonright \xi$$

for all $\xi < \alpha$ and $x, y \in \mathfrak{C}^\alpha$. Put $\mathbb{P}_\alpha = \{\text{range}(f) : f \in \Phi_{\text{prism}}(\alpha)\}$ and $\mathbb{P}_{\omega_1} = \bigcup_{0 < \alpha < \omega_1} \mathbb{P}_\alpha$. The elements of \mathbb{P}_{ω_1} are called the *iterated perfect sets*. If \mathfrak{X} is a Polish space, then a *prism* on X is a pair $\langle f, P \rangle$ where $f: E \rightarrow \mathfrak{X}$ is injective and continuous, $E \in \mathbb{P}_{\omega_1}$, and $P = f[E]$.

It is also immediate to observe that if the pair $\langle f, P \rangle$ is a prism, and $f: E \rightarrow P$ and $E \in \mathbb{P}_\alpha$, then we can assume that f is defined on the entire \mathfrak{C}^α .

It is important to note that the previous definitions imply that perfect cubes are, in particular, iterated perfect sets and therefore, that cubes are prisms. On the other hand, if $\langle g, P \rangle$ is a prism, where $g: E \rightarrow P$ and $E \in \mathbb{P}_\alpha$, then there exists an $f \in \Phi_{\text{prism}}(\alpha)$ with $E = \text{range}(f)$. In particular, $h = g \circ f: \mathfrak{C}^\alpha \rightarrow P$ is a continuous injection and the pair $\langle h, P \rangle$ is a cube. Thus, any prism can be thought of as a cube with a different coordinate system imposed on it.

2.2 Subcubes and Subprisms

If $\langle f, P \rangle$ is a cube, then we say that Q is its subcube provided there exists a perfect cube $C \subset \text{dom}(f)$ such that $Q = f[C]$. Subprisms are defined similarly but with replacing the perfect cube C by an iterated perfect set E . Since in the games defined below we will need to consider singletons in the same position as cubes (or prisms) as defined above, in what follows, *singletons will be considered as cubes and prisms*. If P is a singleton in \mathfrak{X} , then its only subcube is P itself.

2.3 Games and Strategies

For a Polish space \mathfrak{X} , consider the following game $\text{GAME}_{\text{cube}}(\mathfrak{X})$ of length ω_1 played by two players, Player I and Player II. At each

stage $\xi < \omega_1$ of the game Player I can play an arbitrary cube P_ξ in \mathfrak{X} (i.e., P_ξ either belongs to $\text{Perf}(\mathfrak{X})$ or is a singleton in \mathfrak{X}), and Player II must respond by playing a subcube Q_ξ of P_ξ . The game $\langle\langle P_\xi, Q_\xi \rangle : \xi < \omega_1 \rangle$ is won by Player I provided that

$$\mathfrak{X} = \bigcup_{\xi < \omega_1} Q_\xi ;$$

otherwise, Player II wins.

A strategy for Player II is any function S with the property that $S(\langle\langle P_\eta, Q_\eta \rangle : \eta < \xi \rangle, P_\xi)$ is a subcube of P_ξ for every partial game $\langle\langle P_\eta, Q_\eta \rangle : \eta < \xi \rangle$. We say that a game $\langle\langle P_\xi, Q_\xi \rangle : \xi < \omega_1 \rangle$ is played according to a strategy S for Player II provided that $Q_\xi = S(\langle\langle P_\eta, Q_\eta \rangle : \eta < \xi \rangle, P_\xi)$ for every $\xi < \omega_1$. A strategy S for Player II is a *winning strategy* provided Player II wins any game played according to the strategy S . The corresponding notions of games, strategies, etc., for prisms are defined in a similar way.

2.4 The Axioms

The following principles capture the combinatorial core of the iterated Sacks model.

$\text{CPA}_{\text{cube}}^{\text{game}}$: $\mathfrak{c} = \omega_2$ and for any Polish space \mathfrak{X} , Player II has no winning strategy in the game $\text{GAME}_{\text{cube}}(\mathfrak{X})$.

$\text{CPA}_{\text{prism}}^{\text{game}}$: $\mathfrak{c} = \omega_2$ and for any Polish space \mathfrak{X} , Player II has no winning strategy in the game $\text{GAME}_{\text{prism}}(\mathfrak{X})$.

These axioms are consequences of a more general principle, similar in spirit, called CPA [3]. Their importance comes from the following result.

Proposition 2 (Ciesielski and Pawlikowski, [2, 3]). *CPA holds in the iterated perfect set model. In particular, CPA is consistent with ZFC set theory.*

3. AN ω_1 -GENERATED CROWDED Q -POINT ON \mathbb{Q}

If the set $X = [\omega]^{<\omega} \setminus \{\emptyset\}$ has the discrete topology, then the product space $\mathfrak{X} = X^\omega$ is a Polish space, and the family of sets $U_{\langle n, a \rangle} = \{x \in \mathfrak{X} : x(n) = a\}$, where $a \in [\omega]^{<\omega}$ and $n < \omega$, constitutes a subbasis for the product topology. Consider that the set

$$\mathcal{P} = \{x \in \mathfrak{X} : \{x(k) : k < \omega\} \text{ is a partition of } \omega\}.$$

It is important to know that

- \mathcal{P} is a G_δ subset of \mathfrak{X} . Therefore, \mathcal{P} is a Polish space with the relative topology inherited from \mathfrak{X} .

Lemma 1. *Let P be a prism in \mathcal{P} and let $\{A_n : n < \omega\} \subset [\mathbb{Q}]^\omega$ be arbitrary. Then, there exist a subprism Q of P and $B \in [\mathbb{Q}]^\omega$ such that $|B \cap A_n| = \omega$ for every $n < \omega$, and $|x(k) \cap B| \leq 1$ for every $x \in Q$ and $k < \omega$. Moreover, if P is a cube, then Q is a cube as well.*

Proof: Since $|\mathbb{Q}| = \omega$ we can suppose that $\{A_n : n < \omega\} \subset [\omega]^\omega$. Let $\langle R_n : n < \omega \rangle$ be an enumeration of $\{A_n : n < \omega\}$ where each set appears infinitely often.

Case (a): If $P = \{z\}$, then define a sequence $\langle b_n \in \omega : n < \omega \rangle$ such that $b_n \in R_n \setminus \bigcup \{z(k) : k < \omega \text{ \& } z(k) \cap \{b_0, \dots, b_{n-1}\} \neq \emptyset\}$ for every $n < \omega$. It is easy to see that $B = \{b_n : n < \omega\}$ works.

Case (b): If $P \in \text{Perf}(\mathcal{P})$, let f be a witness function for P . By our remarks in section 2, we can assume that f acts from \mathfrak{C}^α onto P . Thus, P is a cube. It is enough to find its subcube with the desired properties.

Let μ be the standard product probability measure on \mathfrak{C}^α .

We construct, by induction on $n < \omega$, a sequence $\langle K_n : n < \omega \rangle$ of open subsets of \mathfrak{C}^α and two sequences, $\langle b_n \in R_n : n < \omega \rangle$ and $\langle B_n \in [\omega]^{<\omega} : n < \omega \rangle$, such that for every $n < \omega$:

- (i) $b_n > \max \left(\{b_i : i < n\} \cup \bigcup_{j < n} B_j \right)$,
- (ii) $\mu(K_n) \geq 1 - 2^{-(n+2)}$, and
- (iii) $f(h)(k) \subseteq B_n$ for every $h \in K_n, k < \omega$ with $b_n \in f(h)(k)$.

If this construction is possible, put $B = \{b_n : n < \omega\}$. Then, clearly $|B \cap A_n| = \omega$. Condition (ii) implies that $\mu \left(\bigcap_{n < \omega} K_n \right) \geq \frac{1}{2}$. Hence, by Proposition 1, there exists a perfect cube $C \subseteq \bigcap_{n < \omega} K_n$. Then $Q = f[C]$ is a subcube of P and the pair $\langle Q, B \rangle$ is as required. To see this, it is enough to show that $|z(k) \cap B| \leq 1$ for every $z \in Q$ and $k < \omega$. Let $z = f(h)$ for some $h \in C$. By conditions (i) and (iii), for every $b_j \in z(k) = f(h)(k)$ and $n > j$, we have that $b_n \notin z(k)$. Therefore, no two elements of B are in the same $z(k)$ or, in other words, $|z(k) \cap B| \leq 1$ for every $k < \omega$.

Next, we show that the inductive construction is possible. Let $n < \omega$ be such that the appropriate b_i , K_i , and B_i are already constructed for every $i < n$. We will construct b_n , K_n , and B_n satisfying (i)–(iii). We pick b_n as an arbitrary element of R_n , satisfying (i). If $L = \{a \in [\omega]^{<\omega} : b_n \in a\}$, then $\{f^{-1}(U_{\langle m, a \rangle}) : \langle m, a \rangle \in \omega \times L\}$ is a partition of \mathfrak{C}^α into clopen sets. Thus, we can find a finite set $S \subseteq \omega \times L$ such that $K_n = \bigcup \{f^{-1}(U_{\langle m, a \rangle}) : \langle m, a \rangle \in S\}$ satisfies condition (ii). Let $B_n = \bigcup \{a : \langle m, a \rangle \in S \text{ for some } m < \omega\}$. Then clearly, B_n is finite. To see that it satisfies (iii), take an $h \in K_n$. Then $f(h) \in U_{\langle m, a \rangle}$ for some $\langle m, a \rangle \in S$. Let $k < \omega$ be such that $b_n \in f(h)(k)$. Since we have also $b_n \in a = f(h)(m)$, we conclude that $k = m$. So, $f(h)(k) = f(h)(m) = a \subseteq B_n$. \square

Fix a $p \in \mathbb{R} \setminus \mathbb{Q}$. For $\mathcal{D} \subset [\mathbb{Q}]^\omega$, let $F(\mathcal{D}) = F(p, \mathcal{D})$ be the filter generated by $\mathcal{D} \cup \{I_n : n < \omega\}$, where $I_n = [p - 2^{-n}, p + 2^{-n}] \cap \mathbb{Q}$.

Lemma 2 (Ciesielski and Pawlikowski [2, Lemma 4.23]). *Suppose that $\mathcal{D} \subset \text{Perf}(\mathbb{Q})$ is a countable family such that $F(\mathcal{D})$ is crowded. Then, for every prism P in $[\mathbb{Q}]^\omega$ there exists a subprism Q of P and a $Z \in \text{Perf}(\mathbb{Q})$ such that $F(\mathcal{D} \cup \{Z\})$ is crowded and either*

- (i) $Z \cap x = \emptyset$ for every $x \in Q$, or else
- (ii) $Z \subset x$ for every $x \in Q$.

We will need also the following easy fact.

Lemma 3 (Ciesielski and Pawlikowski [2, Fact 4.21]). *Every non-scattered set $B \subset \mathbb{Q}$ contains a subset from $\text{Perf}(\mathbb{Q})$.*

Lemma 4. *Let $\mathcal{D} \subset \text{Perf}(\mathbb{Q})$ be a countable family such that $F(\mathcal{D})$ is crowded and let P be a prism in \mathcal{P} , then there exists a subprism Q of P and $Z \in \text{Perf}(\mathbb{Q})$ such that $F(\mathcal{D} \cup \{Z\})$ is crowded and $|Z \cap x(k)| \leq 1$ for every $x \in Q$.*

Proof: Observe that since $F(\mathcal{D})$ is crowded, it is possible to find a sequence $\langle D_n \in \text{Perf}(\mathbb{Q}) : n < \omega \rangle$ coinital in $F(\mathcal{D})$ such that $D_{n+1} \subset D_n \subset I_n$ for every $n < \omega$.

CLAIM. There are sequences $\langle J_k : k < \omega \rangle$ of pairwise disjoint intervals in \mathbb{Q} and $\langle S_k \subset J_k : k < \omega \rangle$ of perfect subsets of \mathbb{Q} such that if $S = \bigcup_{k < \omega} S_k$ then for every $D \in F(\mathcal{D})$ there exists an $n < \omega$ such that $S \cap I_n \subset D$.

To see it, define sequences $\langle n_k : k < \omega \rangle$ and $\langle S_k \in \text{Perf}(\mathbb{Q}) : k < \omega \rangle$ such that $S_k \subset D_k \cap I_{n_k} \cap J_k$ where J_k is a clopen interval such that

$p \notin \text{cl}_{\mathbb{R}}(J_k)$. If n_k and S_k are already defined pick $n_{k+1} > n_k$ with $J_k \cap I_{n_{k+1}} = \emptyset$. Since $D_{k+1} \cap I_{n_{k+1}} \in F(\mathcal{D})$ and $F(\mathcal{D})$ is crowded, we can find a clopen interval J_{k+1} such that $p \notin \text{cl}_{\mathbb{R}}(J_{k+1})$ and $J_{k+1} \cap D_{k+1} \cap I_{n_{k+1}} \neq \emptyset$. Define $S_{k+1} = J_{k+1} \cap D_{k+1} \cap I_{n_{k+1}}$. Then, $S_{k+1} \in \text{Perf}(\mathbb{Q})$ and $S_{k+1} \subset D_{k+1} \cap I_{n_{k+1}}$. Now, put $S = \bigcup_{k < \omega} S_k$. Then, $S \in \text{Perf}(\mathbb{Q})$ and $S \cap I_{n_k} = \bigcup_{i \geq k} S_i \cap I_{n_k} = \bigcup_{i \geq k} S_i \subset D_k$. This proves our claim.

Let \mathcal{B} be a countable basis for the topology on \mathbb{Q} consisting of clopen sets and consider the family $\mathcal{B}_0 = \{B \in \mathcal{B} : |B \cap S| = \omega\}$.

If $P \in \text{Perf}(\mathcal{P})$, apply Lemma 1 to P and $\{B \cap S : B \in \mathcal{B}_0\}$ to find a set $T \in [S]^\omega$ and a subprism Q of P such that

- (a) $|T \cap (B \cap S)| = \omega$ for every $B \in \mathcal{B}_0$ and
- (b) $|T \cap x(k)| \leq 1$ for every $x \in Q$ and $k \in \omega$.

If $P = \{x\}$ is a singleton, we put $Q = P$ and apply Lemma 1 to the family $\{B \cap S : B \in \mathcal{B}_0\}$ and to x to obtain a T satisfying (a) and (b).

In both cases we obtain from (a) that T is dense in S . Since $S_k \in \text{Perf}(\mathbb{Q})$ for every $n < \omega$, we conclude that $T \cap S_k$ is non-scattered and contains a subset Z_k from $\text{Perf}(\mathbb{Q})$ for every $k < \omega$. Hence, if we put $Z = \bigcup_{k < \omega} Z_k$, then $Z \in \text{Perf}(\mathbb{Q})$, $Z \cap I_k \subset D_k$ for every $k < \omega$, and $|Z \cap x(k)| \leq 1$ for every $x \in Q$ and every $k < \omega$. To see that $F(\mathcal{D} \cup \{Z\})$ is crowded, note that $Z \cap D_{n_k} \subset S \cap I_{n_k} \subset D_k$ for every $k < \omega$. \square

Theorem 3. $\text{CPA}_{\text{prism}}^{\text{game}}$ implies that there exists an ω_1 -generated crowded Q -point on \mathbb{Q} .

Proof: For $\mathcal{Y} = [\mathbb{Q}]^\omega \cup \mathcal{P}$, consider the topology τ on \mathcal{Y} whose open sets are those $U \subset \mathcal{Y}$ such that $U \cap [\mathbb{Q}]^\omega$ and $U \cap \mathcal{P}$ are open in $[\mathbb{Q}]^\omega$ and \mathcal{P} , respectively. Then $\langle \mathcal{Y}, \tau \rangle$ is a Polish space. Note that $[\mathbb{Q}]^\omega$ and \mathcal{P} are clopen in \mathcal{Y} with this topology. Every prism $P \in \text{Perf}(\mathcal{Y})$ must intersect either $[\mathbb{Q}]^\omega$ or \mathcal{P} . Since every non-empty clopen set in a prism is its subprism (see [3], or use Proposition 1), we can suppose, without any loss of generality, that either $P \in \text{Perf}([\mathbb{Q}]^\omega)$ or $P \in \text{Perf}(\mathcal{P})$. Of course, every singleton is in either $[\mathbb{Q}]^\omega$ or \mathcal{P} . Therefore, given a prism P in \mathcal{Y} and a countable family $\mathcal{D} \subset \text{Perf}(\mathbb{Q})$ such that $F(\mathcal{D})$ is crowded, we define $Z(\mathcal{D}, P) \in \text{Perf}(\mathbb{Q})$ and a subprism $Q(\mathcal{D}, P)$ of P either as in Lemma 4 if $P \subset [\mathbb{Q}]^\omega$ or as in Lemma 2 if $P \subset \mathcal{P}$.

Consider the following strategy S for Player II:

$$S(\langle\langle P_\eta, Q_\eta \rangle\rangle: \eta < \xi), P_\xi = Q(Z(\{Z_\eta: \eta < \xi\}), P_\xi),$$

where sets Z_η are defined inductively by $Z_\eta = Z(\{Z_\zeta: \zeta < \eta\}, P_\eta)$.

By CPA_{prism}^{game} strategy, S is not a winning strategy for Player II. Hence, there is a game $\langle\langle P_\xi, Q_\xi \rangle\rangle: \xi < \omega_1$ played according to S for which Player II loses; so $\mathcal{Y} = \bigcup_{\xi < \omega_1} Q_\xi$.

Let $\mathcal{U} = F(\{Z_\xi: \xi < \omega_1\})$. To see it is an ultrafilter, note that if $x \in [\mathbb{Q}]^\omega$ then there exists a $\xi < \omega_1$ such that $x \in Q_\xi$. But then, either $Z_\xi \subset x$ or $Z_\xi \cap x = \emptyset$. Therefore, either x or its complement is in \mathcal{U} . This proves that \mathcal{U} is an ultrafilter and that $\langle Z_\xi: \xi < \omega_1 \rangle \subset \text{Perf}(\mathbb{Q})$ is a basis for \mathcal{U} . Therefore, \mathcal{U} is crowded. Since no crowded ultrafilter can be principal, it follows that \mathcal{U} is also non-principal. To see that \mathcal{U} is a Q -point, pick an $x \in \mathcal{P}$. Then there exists a $\xi < \omega_1$ such that $x \in Q_\xi$. Thus, $Z_\xi \in \mathcal{U}$ and $|Z_\xi \cap x(k)| \leq 1$ for every $k < \omega$. \square

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