

Counting Weakly Bi-colored Trees

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Outline

1 Coloring

2 Connections

3 Counts

Outline

1 Coloring

2 Connections

3 Counts

Outline

- 1 Coloring
- 2 Connections
- 3 Counts

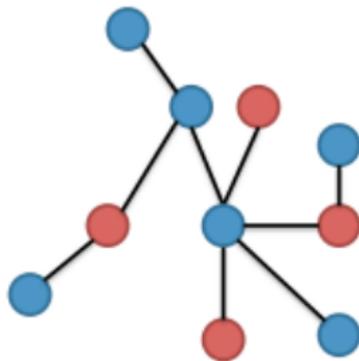
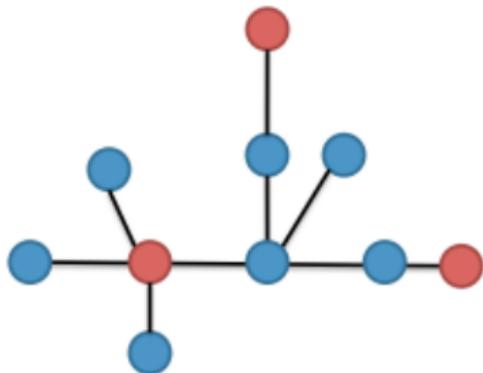
Weakly Bi-colored Trees

Definition

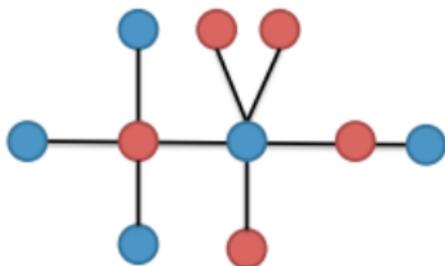
A tree is said to be *weakly bi-colored* provided it satisfies the following conditions:

- 1 Two colors (say, red and blue) that both must be used at least once.
- 2 Adjacent vertices can both be blue.
- 3 Adjacent vertices cannot both be red.

Examples of Weakly Bi-colored Trees



Examples of Weakly Bi-colored Trees



- A strongly bi-colored tree, which is considered weakly bi-colored

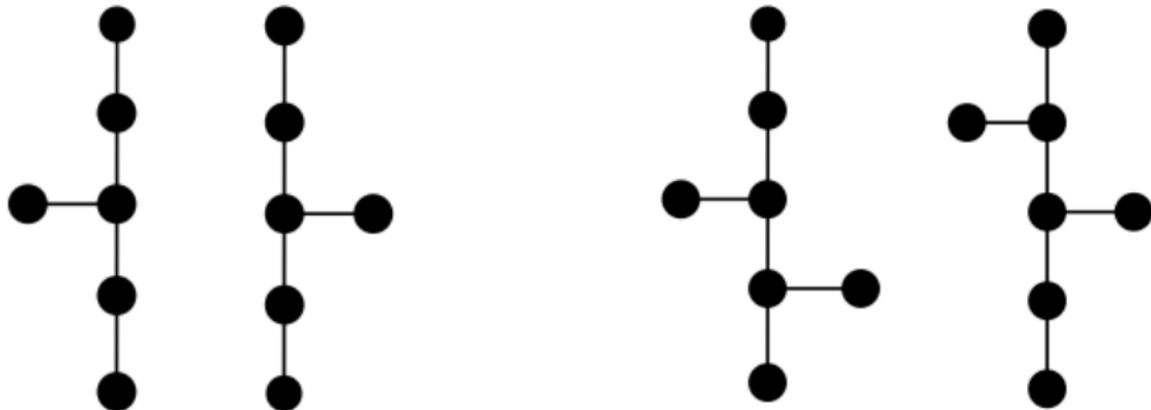
Planar Equivalent Trees

Definition

Two trees are said to be *planar equivalent* (or *planar isomorphic*) if and only if there is an orientation-preserving homeomorphism of the plane between them.

Problem: This condition complicates the count.

Examples of Planar Equivalent Trees



Color Equivalent Trees

Definition

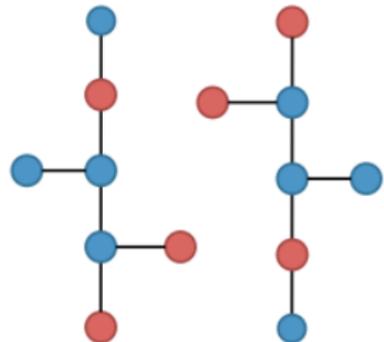
Two trees are said to be *color equivalent* (or *color isomorphic*) if and only if the mapping between the trees respects the coloring.

Definition

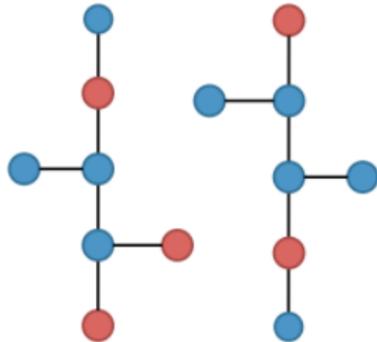
If two trees are both planar equivalent and color equivalent, we will say that the trees are *equivalent* or *isomorphic*.

- From now on, we will use the terms equivalent and isomorphic interchangeably

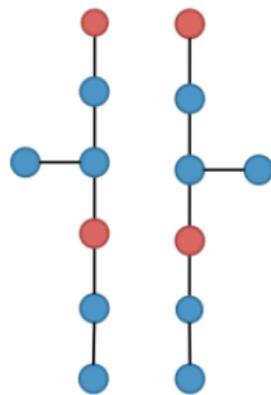
Examples of Color Planar Equivalent Trees



Color, Planar



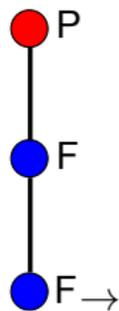
Non-color, Planar



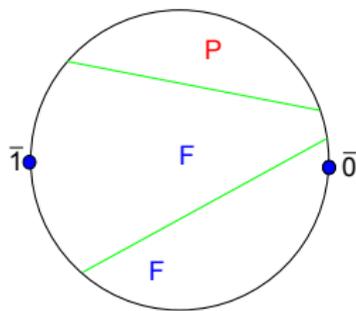
Color, Non-planar

Levels of Correspondence

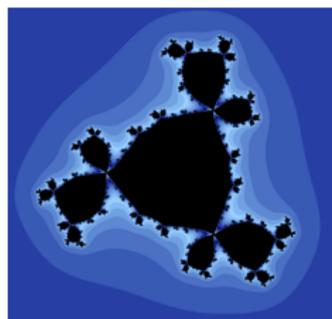
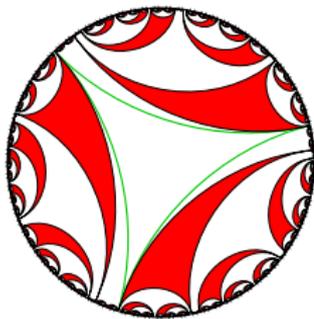
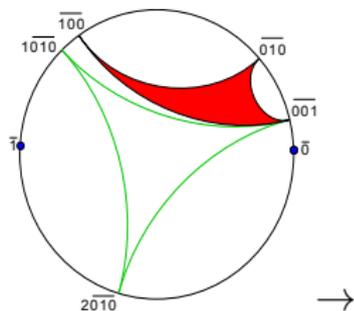
Bi-colored Tree



Critical Portrait



Lamination Data



Lamination

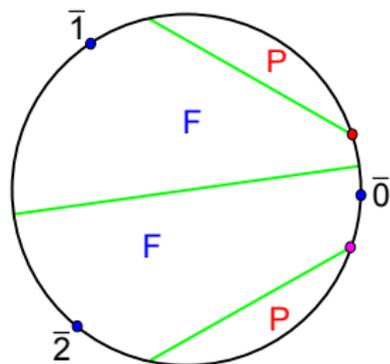
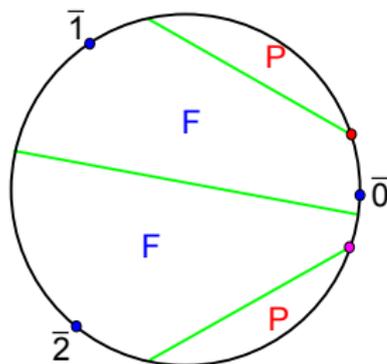
Julia Set

Correspondence Between Bi-colored trees and Critical Portraits

- The correspondence between weakly bi-colored trees and critical portraits raises some interesting questions
- Going from a critical portrait to a tree is unambiguous, there is only one weakly bi-colored tree for a given critical portrait.
- However, going from a weakly bi-colored tree to critical portrait there is ambiguity. A tree may correspond to multiple portraits.

Different Critical Portraits with Same Bi-colored Tree

The two critical portraits below correspond to the same weakly bicolored tree:



Strongly Bi-colored Case

Definition

A tree is said to be *strongly bi-colored* provided it satisfies the following conditions:

- 1 Two colors (say, red and blue) that both must be used at least once.
- 2 Adjacent vertices cannot both be the same color.

Strongly Bi-colored Case

- The number of strongly bi-colored planar trees with d edges is:

$$N(d) = \frac{1}{d} \left(\frac{1}{d+1} \binom{2d}{d} + \sum_{n|d, n < d} \phi\left(\frac{d}{n}\right) \binom{2n}{n} \right) \quad [1] \quad (1)$$

- Where $\phi(x)$ is Euler's totient function, the number of positive integers less than, and relatively prime to, x .

[1]-Casper et al, 2015

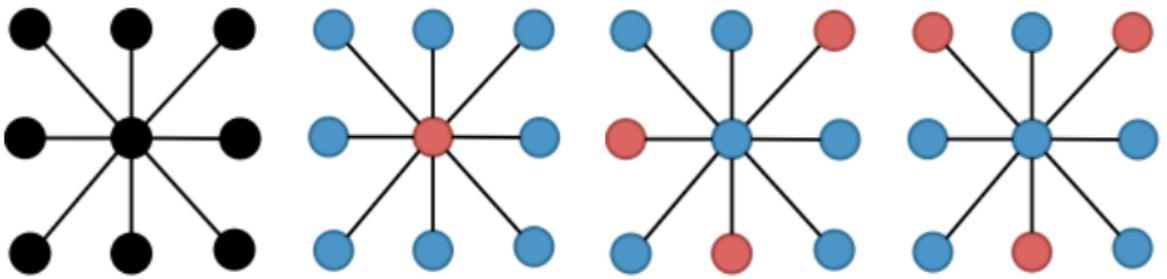
Fans

Definition

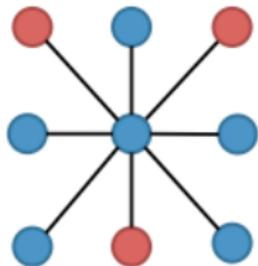
To be classified as a *fan* a tree with $n + 1$ vertices must have the following qualities:

- One central vertex of degree n
- n vertices with degree 1

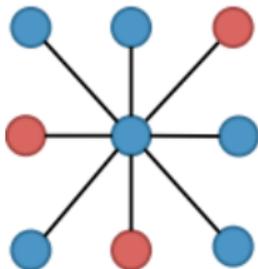
Examples of Fans



Representing Fans



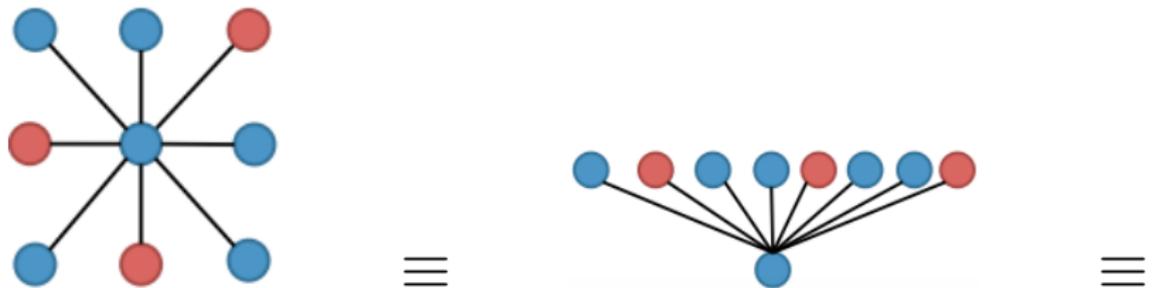
≡



≡



Representing Fans



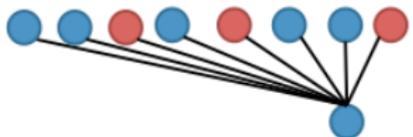
Representing Fans

- We can represent Fans as a binary sequence with 1's representing reds and 0's representing blues



≡

01001001



≡

00101001

Representing Fans

- We can take all rotations to be equivalent.

$$01001001 \equiv 00101001 \equiv 10010100 \equiv 01001010 \equiv \\ 00100101 \equiv 10010010 \equiv 01001001$$

- Taking all rotations to be a equivalent means that we can reduce *Fans* to the combinatorial objects, *Necklaces*

Representing Fans

Definition

A k -ary *necklace* is an array of length n with k distinct characters taking all rotations to be equivalent.[2]

- The number of necklaces with a given n and k is:

$$N_k(n) = \frac{1}{n} \sum_{d|n} \phi(d) k^{\frac{n}{d}} \quad [2] \quad (2)$$

- Where $\phi(x)$ is Euler's totient function, the number of positive integers less than, and relatively prime to, x .

[2]-Weisstein

Representing Fans

- A fan with $n + 1$ vertices can be represented as a necklace of length n with $k = 2$ distinct characters (1 and 0)
- The number of weakly bi-colored fans with $n + 1$ vertices (or n edges) is:

$$N(n) = \frac{1}{n} \sum_{d|n} \phi(d) 2^{\frac{n}{d}} \quad (3)$$

Chains

Definition

A *Chain* is a tree with the following qualities:

- two terminal vertices of degree one
- Any non-terminal vertex is of degree two

Examples of Chains



Representing Chains

- Chains can also be represented with binary sequences of 1's and 0's.
- 0's representing blues, 1's representing reds and 1's cannot be sequential



Counting Chains

- We can think of the problem of counting chains by trying to count to the number of sequences of 1's and 0's where 1's are not allowed to touch

Counting Chains

- Say we have n vertices (or slots in the sequence)
- Let i be the number of reds (or 1's) with $1 \leq i \leq \lceil n/2 \rceil$
- Let $n - i$ be the number of blues (or 0's)

Counting Chains

- Now suppose we were to put down all of the 0's (or blues):
_0_0_0_..._0_0_
- This gives us $n - i + 1$ slots to place the i 1's (or reds) and once there is a 1 (or red) in a slot nothing else can be placed in said slot
- Meaning we have $n - i + 1$ places and i objects to place, so there are $\binom{n-i+1}{i}$ ways to place the 1's (or reds)

Counting Chains

- However $\binom{n-i+1}{i}$ will overcount because rotation is not considered in this number.
- Some trees will be counted twice.
- Some trees will be counted once and these will be called *palindromes*



Palindromes and Non-palindromes

- Non-palindromes are counted twice and palindromes are counted once due to the condition of planar equivalence.

Non-Palindromes



Palindrome



Counting Palindromes

- In order to correct the overcounting, we need to count the palindromes.
- We can do this in a similar process as how we found the overall count.
- This time we have to consider only the first half of the chain and whether the number of vertices and the number of reds are even or odd.

Counting Palindromes

- We place down half of the blues and place half of the reds in between blues, then flip it across horizontally to create a palindrome.
- So, by counting the ways of arranging the first half, we can count the number of palindromes.



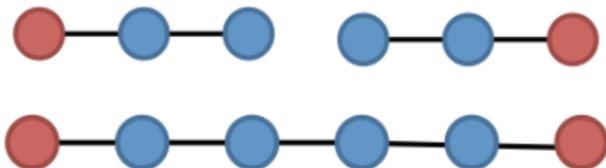
Counting Palindromes

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Counting Palindromes

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- So, by counting the ways of arranging the first half, we can count the number of palindromes.



Counting Palindromes

- 1 If n is even and i is odd, there will be no palindromes.
- Recall $n =$ number of vertices and $i =$ number of reds

Counting Palindromes

- 1 If n is even and i is odd, there will be no palindromes.
 - 2 If n is even and i is even, there will be $\binom{\frac{n-i}{2}}{\frac{i}{2}}$ palindromes.
- Recall $n =$ number of vertices and $i =$ number of reds

Counting Palindromes

- 1 If n is even and i is odd, there will be no palindromes.
 - 2 If n is even and i is even, there will be $\binom{\frac{n-i}{2}}{\frac{i}{2}}$ palindromes.
 - 3 If n is odd and i is odd, there will be $\binom{\frac{n-i}{2}}{\frac{i-1}{2}}$ palindromes.
- Recall $n =$ number of vertices and $i =$ number of reds

Counting Palindromes

- 1 If n is even and i is odd, there will be no palindromes.
 - 2 If n is even and i is even, there will be $\binom{\frac{n-i}{2}}{\frac{i}{2}}$ palindromes.
 - 3 If n is odd and n is odd, there will be $\binom{\frac{n-i}{2}}{\frac{i-1}{2}}$ palindromes.
 - 4 If n is odd and i is even, there will be $\binom{\frac{n-i+1}{2}}{\frac{i}{2}}$ palindromes.
- Recall $n =$ number of vertices and $i =$ number of reds

Final Count for Chains

- By combining the two counts we found, we have that the number of weakly bi-colored planar chains with n vertices is:

$$N(n) = \frac{1}{2} \sum_{i=1}^{\lceil \frac{n}{2} \rceil} \left(\binom{n-i+1}{i} + P(n, i) \right) \quad (4)$$

Where $P(n, i) =$ number of palindromes:

$$P(n, i) = \begin{cases} 0, & n \text{ even and } i \text{ odd} \\ \binom{\frac{n-i}{2}}{i/2}, & n \text{ even and } i \text{ even} \\ \binom{\frac{n-i}{2}}{(i-1)/2}, & n \text{ odd and } i \text{ odd} \\ \binom{\frac{n-i+1}{2}}{i/2}, & n \text{ odd and } i \text{ even} \end{cases} \quad (5)$$

Caterpillars and Lobsters

Definition

A *Caterpillar* is a tree with the following qualities:

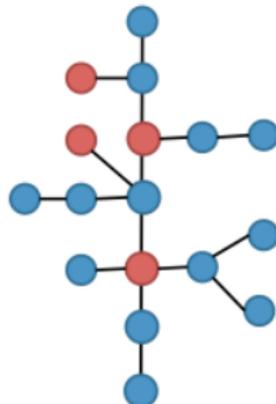
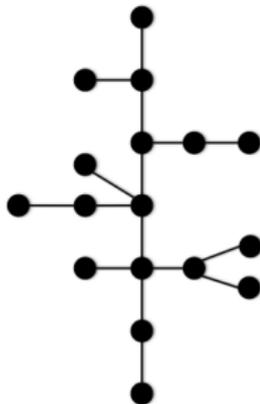
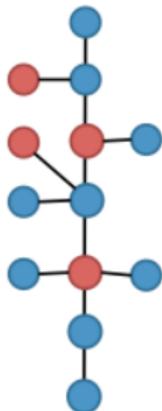
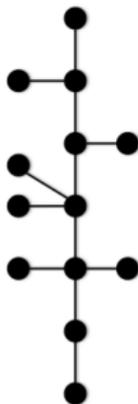
- There is some longest backbone chain of length k ,
 $4 \leq k \leq n - 1$
- Any vertex not on the backbone is a distance of one off the backbone and does not increase the backbone length.

Definition

A *Lobster* is a tree with the following qualities:

- There is some longest backbone chain of length k ,
 $4 \leq k \leq n - 1$
- Any vertex not on the backbone is a distance of at most two off the backbone and does not increase the backbone length.

Examples of Caterpillars and Lobsters

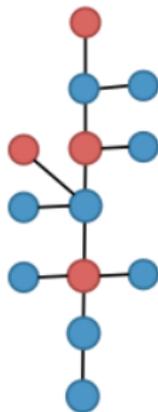
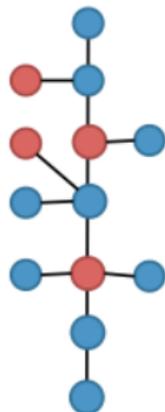


Problems arising from Caterpillars and Lobsters

- In a caterpillar or lobster, the longest backbone is not necessarily unique.
- Two trees can be equivalent and have backbones of the same length, but the backbones are different colors.
- This is something that needs to be considered when making the count.

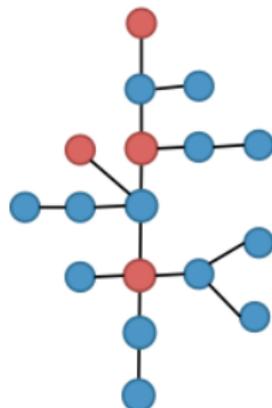
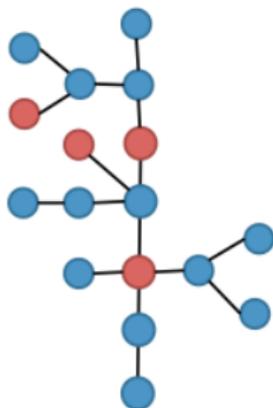
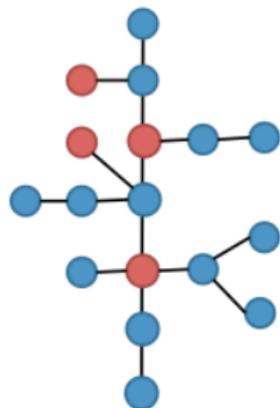
Problems Arising from Caterpillars and Lobsters

- Two caterpillars that are equivalent, but have different backbones.



Problems Arising from Caterpillars and Lobsters

- Three equivalent lobsters, each with a different backbone.



Super-Lobster

Definition

A *Super-Lobster* is a tree with the following qualities:

- There is some longest backbone chain of length k ,
 $3 \leq k \leq n - 1$
 - Any vertex not on the backbone is a length of m ,
 $1 \leq m \leq \lfloor \frac{k}{2} \rfloor$, off the backbone and does not increase the backbone length.
-
- Any tree is representable as a *super-lobster*
 - Same problem as caterpillars and lobsters, the longest backbone may not be unique.

Caterpillars, Lobster, and Super-Lobsters

- Currently, the count for caterpillars is being explored.
- Count for caterpillars \rightarrow count for lobsters \rightarrow count for super-lobsters

